Convex Optimization

(EE227A: UC Berkeley)

Lecture 20 (Coordinate descent)

04 Apr, 2013

Suvrit Sra

Admin

- \heartsuit HW3 due right now!
- \heartsuit HW4 is out! Please ask your Qs on Piazza
- \heartsuit Project 4 page reports due on 4/11/2013
- \heartsuit Poster presentations: 3hrs: When in May?

Challenge problem

$$I(p) := \sqrt{p} \int_0^\infty \left| \frac{\sin x}{x} \right|^p dx$$

Minimize I(p) over $p \ge 1$

So far:
$$\min f(x) = \sum_i f_i(x)$$

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Since $x \in \mathbb{R}^n$, now consider $\min f(x) = f(x_1, x_2, \dots, x_n)$

Previously, we went through f_1,\ldots,f_m

What if we now go through x_1, \ldots, x_n one by one?

For
$$k = 0, 1, ...$$

Coordinate descent

$$\blacksquare \quad \text{For } k = 0, 1, \dots$$

Pick an index i from $\{1, \ldots, n\}$

Coordinate descent

For
$$k = 0, 1, ...$$

Pick an index i from $\{1, ..., n\}$
Optimize the i th coordinate
 $x_i^{k+1} \leftarrow \operatorname*{argmin}_{\xi \in \mathbb{R}} f(\underbrace{x_1^{k+1}, \ldots, x_{i-1}^{k+1}}_{\text{done}}, \underbrace{\xi}_{\text{current}}, \underbrace{x_{i+1}^k, \ldots, x_n^k}_{\text{todo}})$

Decide when/how to stop; return x^k

 x_i^{k+1} overwrites value in x_i^k (in actual implementation)

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- \clubsuit These days renewed interest in CD for large-scale problems
- Notice: in general CD is "derivative free"

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Derivative free rules:

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- **Random sampling**: pick random index at each iteration

Coordinate descent – Example

$$\min \|Ax - b\|_2^2$$

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Coordinate descent update

$$x_j \leftarrow \frac{\sum_{i=1}^m a_{ij} \left(b_i - \sum_{l \neq j} a_{il} x_l \right)}{\sum_{i=1}^m a_{ij}^2}$$

(dropped superscripts, since we overwrite)

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- Nonsmooth case more tricky
Block coordinate descent (BCD)

$$\begin{array}{ll} \min & f(\boldsymbol{x}) := f(\boldsymbol{x}_1, \dots, \boldsymbol{x}_m) \\ & \boldsymbol{x} \in \mathcal{X}_1 \times \mathcal{X}_2 \times \dots \times \mathcal{X}_m. \end{array}$$

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Gauss-Seidel style



Jacobi style (easy to parallelize)

$$\boldsymbol{x}_{i}^{k+1} \leftarrow \operatorname*{argmin}_{\boldsymbol{\xi} \in \mathcal{X}_{i}} f(\underbrace{\boldsymbol{x}_{1}^{k}, \ldots, \boldsymbol{x}_{i-1}^{k}}_{\text{don't clobber}}, \underbrace{\boldsymbol{\xi}}_{\text{current}}, \underbrace{\boldsymbol{x}_{i+1}^{k}, \ldots, \boldsymbol{x}_{m}^{k}}_{\text{todo}})$$

BCD – convergence

Theorem Let f be continuously differentiable over $\mathcal{X} := X_{i=1}^m \mathcal{X}_i$. Further, assume for each block i and $x \in \mathcal{X}$, the minimum

$$\min_{\boldsymbol{\xi}\in\mathcal{X}_i} f(\boldsymbol{x}_1,\ldots,\boldsymbol{x}_{i+1},\boldsymbol{\xi},\boldsymbol{x}_{i+1},\ldots,\boldsymbol{x}_m)$$

is **uniquely attained**. Every limit point of the sequence $\{x^k\}$ generated by BCD, is a stationary point of f.

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Corollary. If f is in addition convex, then every limit point of the BCD sequence $\{x^k\}$ is a global minimum.

- ► Unique solutions of subproblems not always possible
- Above result is only **asymptotic** (holds in the limit)
- ▶ Warning! BCD may cycle indefinitely without converging, if the number of blocks is > 2 and the objective is nonconvex.

Two block BCD

minimize $f(\boldsymbol{x}) = f(\boldsymbol{x}_1, \boldsymbol{x}_2) \quad \boldsymbol{x} \in \mathcal{X}_1 \times \mathcal{X}_2.$

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Theorem (Grippo & Sciandrone (2000)). Let f be continuously differentiable, and the sets \mathcal{X}_1 , \mathcal{X}_2 be closed and convex. Assume that the both BCD subproblems have solutions, and that the sequence $\{x^k\}$ has limit points. Then, every limit point of $\{x^k\}$ is stationary.

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- ▶ No need of **unique solutions** to subproblems
- ▶ BCD for 2 blocks is also called: Alternating Minimization

CD for convex problems

CD for smooth convex problems

$\min f(Ax) + \langle b, x \rangle \text{ subject to } x \ge 0$

- \blacktriangleright Function f is strictly convex and smooth
- Matrix $A \in \mathbb{R}^{m \times n}$ (possibly rank-deficient)

CD for smooth convex problems

$\min f(Ax) + \langle b, x \rangle$ subject to $x \ge 0$

- Function f is strictly convex and smooth
- Matrix $A \in \mathbb{R}^{m \times n}$ (possibly rank-deficient)
- ► Apply CD to this problem
- ▶ With some more assumptions: it works!
- ► Even rate of convergence analysis (asymptotic)
- ► Here's the theorem

Assumptions:

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Theorem (Luo, Tseng (1992)). Let $\{x^k\}$ be a sequence of iterates generated by the CD method using the almost cyclic or the Gauss-Southwell rule for picking indices. Then, $\{x^k\}$ converges at least linearly to an element of \mathcal{X}^* .

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Proof is intricate; see Luo & Tseng's paper on bSpace.

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$$\frac{1}{2} \|x - y\|_2^2 + \sum_{i=1}^m \delta_{C_i}(x).$$

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► Now invoke Douglas-Rachford using the product-space trick Solution 2: Take dual of the above formulation

Convex calculus time

$$\min \quad \frac{1}{2} \|x - y\|_2^2 + f(x) + h(x)$$

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 $L(x,z,w,\nu,\mu) := \frac{1}{2} \|x-y\|_2^2 + f(z) + h(w) + \nu^T (x-z) + \mu^T (x-w)$

$$g(\nu, \mu) := \inf_{\substack{x, z, w}} L(x, z, \nu, \mu)$$

$$x - y + \nu + \mu = 0 \implies x = y - \nu - \mu$$

$$g(\nu, \mu) = -\frac{1}{2} \|\nu + \mu\|_2^2 + (\nu + \mu)^T y - f^*(\nu) - h^*(\mu)$$

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Dual as minimization problem

min
$$k(\nu,\mu) := \frac{1}{2} \|\nu + \mu - y\|_2^2 + f^*(\nu) + h^*(\mu)$$

Apply CD to $k(\nu, \mu)$

 $\begin{array}{lll} \nu_{k+1} &=& \mathrm{argmin}_{\nu} \ k(\nu,\mu_k) \\ \mu_{k+1} &=& \mathrm{argmin}_{\mu} \ k(\nu_{k+1},\mu) \end{array}$

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$$b \quad 0 \in \nu + \mu_k - y + \partial f^*(\nu) b \quad 0 \in \nu_{k+1} + \mu - y + \partial h^*(\mu) b \quad y - \mu_k \in \nu + \partial f^*(\nu) = (I + \partial f^*)(\nu) \implies \nu = \operatorname{prox}_{f^*}(y - \mu_k) \implies \nu = y - \mu_k - \operatorname{prox}_f(y - \mu_k)$$

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► $0 \in \nu + \mu_k - y + \partial f^*(\nu)$ ► $0 \in \nu_{k+1} + \mu - y + \partial h^*(\mu)$ ► $y - \mu_k \in \nu + \partial f^*(\nu) = (I + \partial f^*)(\nu)$ $\implies \nu = \operatorname{prox}_{f^*}(y - \mu_k) \implies \nu = y - \mu_k - \operatorname{prox}_f(y - \mu_k)$ ► Similarly, we see that

$$\mu = y - \nu_{k+1} - \operatorname{prox}_h(y - \nu_{k+1})$$

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Proximal-Dykstra as CD

■ Simplify, and use Lagrangian stationarity to obtain primal

$$x = y - \nu - \mu \implies y - \mu = x + \nu$$

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■ Thus, the CD iteration may be rewritten as

$$t_k \leftarrow \operatorname{prox}_f(x_k + \nu_k)$$
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• We used: $\operatorname{prox}_{h}(y - \nu_{k+1}) = \mu_{k+1} - y - \nu_{k+1} = x_{k+1}$
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We used: prox_h(y − ν_{k+1}) = µ_{k+1} − y − ν_{k+1} = x_{k+1}
This is the proximal-Dykstra method!

CD for nonsmooth convex problems



CD for separable nonsmoothness

► Nonsmooth part is **separable**

$$\min_{x \in \mathbb{R}^n} f(x) + \sum_{i=1}^n r_i(x_i)$$

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Remark: A related result for **nonconvex** problems with separable non-smoothness (under more assumptions), can be found in: "*Convergence of Block Coordinate Descent Method for Nondifferentiable Minimization*" by P. Tseng (2001).

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- Assume *f* is convex, with **componentwise** Lipschitz gradients

$$|\nabla_i f(x+he_i) - \nabla_i f(x)| \le L_i |h|, \quad x \in \mathbb{R}^n, h \in \mathbb{R}.$$

Here e_i denotes the *i*th canonical basis vector

- ► So far, we saw CD based on essentially cyclic rules
- It is difficult to prove global convergence and almost impossible to estimate global rate of convergence
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Choose
$$x_0 \in \mathbb{R}^n$$
. Let $M = \max_i L_i$; For $k \ge 0$
 $i_k = \operatorname{argmax}_{1 \le i \le n} |\nabla_i f(x_k)|$
 $x_{k+1} = x_k - \frac{1}{M} \nabla_{i_k} f(x_k) e_{i_k}$.

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- ▶ Also, if $f \in C^1_L$, it can easily happen that $M \ge L$
- ▶ So above rate is in general, worse than gradient methods

NEXT LECTURE:

- ► Randomized BCD (aka Stochastic BCD)
- ► Parallel BCD
- ► Dual decomposition, ADMM, etc.